

MIPS arithmetic

- Unsigned and signed number representations.
 - Addition and subtraction with two's complement numbers.
 - Overflow detection.
- f* - review the construction of an ALU that will appear in our CPU Design

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Unsigned numbers

- f* We can store unsigned numbers as their binary equivalents.
f Bit position i has a decimal weight 2^i , so the decimal value v of an n -bit unsigned binary number $a_{n-1}a_{n-2}\dots a_1a_0$ can be calculated as below.

$$v = \sum_{i=0}^{n-1} 2^i a_i$$

- f*This is just the sum of each digit times the weight of its position.

$$101101_2 = (2^5 \cdot 1) + (2^4 \cdot 0) + (2^3 \cdot 1) + (2^2 \cdot 1) + (2^1 \cdot 0) + (2^0 \cdot 1) = 45_{10}$$

- f* The smallest and largest possible numbers are 0 and $2^n - 1$.
*f*n MIPS, the largest unsigned number is $2^{32} - 1$, or about 4.3 billion.

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Hexadecimal notation

f Hexadecimal is frequently used as a shorthand for binary numbers, because one hex digit corresponds to four bits.

f Converting between binary and hex is easy with a table.

$$\text{BEEF}_{16} = 1011\ 1110\ 1110\ 1111_2$$

*f*n SPIM and many high-level programming languages, the prefix **0x** denotes a hexadecimal constant, as in **0xBEEF**.

Binary	Hex
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	8
1001	9
1010	A
1011	B
1100	C
1101	D
1110	E
1111	F

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Working with unsigned numbers

f Addition can be done just like in decimal.

$$\begin{array}{r}
 1\ 1\ 1\ 1\ 1 \\
 1\ 0\ 1\ 1\ 0\ 1 \\
 +\ 0\ 1\ 0\ 1\ 1\ 1 \\
 \hline
 1\ 0\ 0\ 0\ 1\ 0\ 0 \\
 \hline
 \end{array}
 \qquad
 \begin{array}{r}
 45 \\
 +\ 23 \\
 \hline
 68
 \end{array}$$

f Multiplication and integer division by powers of 2 can be done using left and right “logical” shifts.

$$\begin{array}{l}
 101101 \times 10 = 1011010 \\
 101101 / 10 = 010110
 \end{array}
 \qquad
 \begin{array}{l}
 45 \times 2 = 90 \\
 45 / 2 = 22
 \end{array}$$

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Logical shifts in MIPS

- f* MIPS has `sll` (shift left logical) and `srl` (shift right logical) instructions.
 - A constant specifies the number of bits to shift, from 0 to 31.
 - 0s are shifted into the right or left sides, respectively.
- f* For example, assume that `$t0` contains `0xFEEDBEEF`.

```
sll $t1, $t0, 4    # $t1 contains 0xEEDBEEF0
srl $t2, $t0, 12   # $t2 contains 0x00FEEDB
```

f Shifts are actually R-type instructions, not I-type!

op	rs	rt	rd	shamt	func
6 bits	5 bits	5 bits	5 bits	5 bits	6 bits

f The constant is stored in the 5-bit “shamt” field that we hadn’t seen up until now.

Unsigned overflow

- f* One recurring issue in computer arithmetic is dealing with finite amounts of storage, such as 32-bit MIPS registers.
 - f* **Overflow** occurs when the result of an operation is too large to be stored.
 - f* There are many examples of unsigned n -bit overflows.
 - When there is a carry out of position $n-1$ in an addition.

f The 6-bit addition and multiplication operations on page 4 both result in overflow, since the correct answers require 7 bits.

Signed two's-complement numbers

Signed numbers are represented in **two's complement** format.

The most significant bit a_{n-1} of each number is a **sign bit**.

- 0 indicates a positive number.
- 1 indicates a negative number.

The **range** of n -bit signed numbers is from -2^{n-1} to $+2^{n-1}-1$

- For 32-bit values, the range is roughly ± 2.15 billion.
- The ranges are not exactly even since there are 2^{n-1} negative numbers, but just $2^{n-1}-1$ positive numbers.

Negating a two's complement number

There are three main ways to negate two's complement numbers. As an example, let's consider the six-bit value **101101** (-19).

1. Complement all the bits in the number and then add 1.
 - Complementing the bits in 101101 yields 010010.
 - Adding 1 results in 010011 (+19).
2. Complement the bits to the left of the rightmost 1.
 - The rightmost 1 in 101101 is in position 0.
 - Complementing the bits to its left gives 010011 (+19).
3. Subtract the number from 2^n .
 - 2^6 is 1000000.
 - Then, $1000000 - 101101 = 010011$ (+19).

Two's-complement numbers with $n=4$

Decimal	2C		2C	Decimal
-8	1000	↔	0000	0
-7	1001		0001	1
-6	1010		0010	2
-5	1011		0011	3
-4	1100		0100	4
-3	1101		0101	5
-2	1110		0110	6
-1	1111		0111	7
0	0000		1000	-8
1	0001		1001	-7
2	0010		1010	-6
3	0011		1011	-5
4	0100		1100	-4
5	0101		1101	-3
6	0110		1110	-2
7	0111		1111	-1

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Working with signed numbers

Addition with two's complement is the same as with unsigned numbers.

- The sign bits are included in the addition.
- The carry out is ignored.

$$\begin{array}{r}
 \\
 \\
 + \\
 \hline
 1
 \end{array}
 \begin{array}{r}
 \\
 \\
 \\
 \\
 + 23 \\
 \hline
 4
 \end{array}$$

A subtraction operation $A - B$ is treated as an addition of $A + (-B)$.

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Sign extension

*f*We often work with signed numbers of different lengths.

- Instructions like `lw $t0, -4($sp)` add 16-bit and 32-bit values together.
- `slti $t0, $t0, 256` compares a 16-bit constant to a 32-bit number. The
- `lb` instruction copies an 8-bit value into a 32-bit register.

f You can **sign extend** a two's complement number by copying the sign bit.

f For instance, we can sign extend 6-bit numbers to 8 bits.

(+23) 010111 → 00010111 (+23)

(-19) 101101 → 11101101 (-19)

*f*f you add 0s instead of sign extending, you may accidentally change the sign, and hence the value, of the result.

(-19) 101101 → 00101101 (+45)



Interpreting bit patterns

f One of the most important things to remember is that bit patterns have different meanings under different representations!

- As a six-bit *unsigned* number, 101101 denotes 45 in decimal.
- But as a *two's complement* number, 101101 denotes -19.

*f*The example C program below prints the same data (0xFFFFFFFF) twice, but under different interpretations.

```
main()                // CSIL Sun machines
{
  int x = 0xFFFFFFFF; // 32-bit integers
  printf("x = %d\n", x); // Signed; prints -1
  printf("x = %u\n", x); // Unsigned; prints 4294967295
}
```

Signed and unsigned comparisons

f Specifying the interpretation of binary numbers is especially important in untyped assembly language programming.

f For example, how does 111111 compare to 000000?

- As an unsigned number, 111111 = 63, which is *greater* than 0.
- As a two's complement number, 111111 = -1, which is *less* than 0.

f MIPS includes two versions of the comparison instructions.

- The `slt` and `slti` instructions that we saw do signed comparisons.
- The instructions `sltu` and `sltiu` perform unsigned comparisons.

```
main()
{
    unsigned int u = 0xFFFFFFFF;
    int s = 0xFFFFFFFF;
    printf("%u is %s than 0\n", u, (u < 0 ? "less" : "greater"));
    printf("%d is %s than 0\n", s, (s < 0 ? "less" : "greater"));
}
```

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Signed and unsigned addition

f MIPS provides two versions of the main arithmetic instructions.

- The `add`, `addi` and `sub` instructions will raise an exception if they cause an overflow.

- But `addu`, `addiu`, and `subu` do not raise exceptions on overflow

f Be careful! Both `addi` and `addiu` sign extend their constant field, even though `addiu` is considered an “unsigned” operation.

f In C, overflow exceptions may or may not be raised, depending on the underlying hardware.

```
main() // CSIL Sun machines
{
    int x = 0x7FFFFFFF; // Largest 32-bit number
    printf("x = %d\n", x); // Prints 2147483647
    x++;
    printf("x = %d\n", x); // Prints -2147483648
}
```

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When to use unsigned numbers

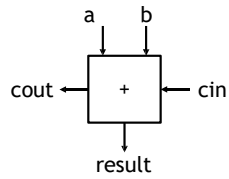
- f* Memory addresses should be treated as unsigned numbers.
 - The fourth element of an array that starts at 0x7FFFFFF0 is at address 0x80000000. This address computation should not overflow!
 - In fact, stack pushes and pops should be performed using the unsigned add and subtract instructions like `addiu`, and not `addi`.
- f* Non-numeric data should also be handled using unsigned operations.
 -

MIPS arithmetic instruction summary

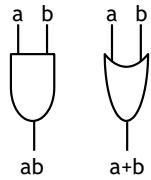
- f* Addition and subtraction
 - `add`, `addi`, and `sub` raise exceptions on overflow.
 - `addu`, `addiu`, and `subu` do not raise exceptions on overflow.
 - `addi` and `addiu` both sign extend their immediate fields.
- f* Comparisons
 - `slt` and `slti` are for signed comparisons.
 - `sltu` and `sltiu` do unsigned comparisons.
 - `slti` and `sltiu` both sign extend their immediate fields.
- f* Data transfers
 - `lb` sign extends the byte that is loaded into a register.
 - `lbu` does not sign extend the byte.
- f* Shifting
 - `sll` and `srl` are used for logical shifts of up to 31 bits.
- f* Bitwise operators
 - `andi` and `ori` do not sign extend their immediate field.

Hardware for single-bit operations

We can use a full adder to add data inputs *a*, *b* and a carry in *cin*. This produces a one-bit *result* and a carry out *cout*.



It's also easy to do one bit *and* and *or* operations.



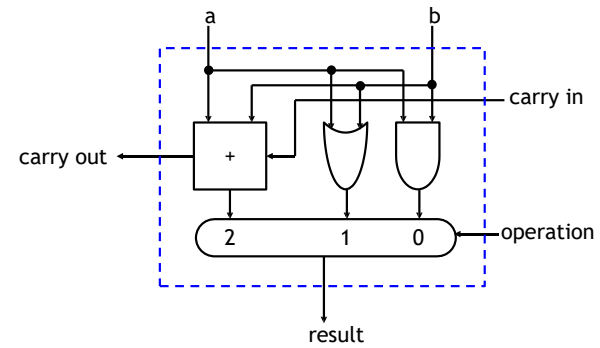
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A one-bit ALU

A simple one-bit ALU can support all three of these functions.

A two-bit input called *operation* selects the desired function using a 3-to-1 multiplexer.

operation	result
0	<i>a</i> and <i>b</i>
1	<i>a</i> or <i>b</i>
2	<i>a</i> + <i>b</i> + carry in

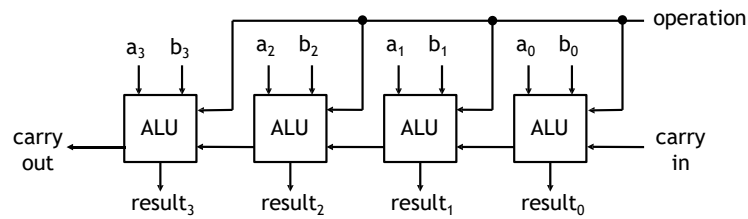


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An n -bit ALU

This one-bit ALU can be replicated to produce an n -bit ALU.

- The **carry out** from stage i connects to the **carry in** of stage $i+1$.
- The **operation** input is shared by all of the one-bit circuits.



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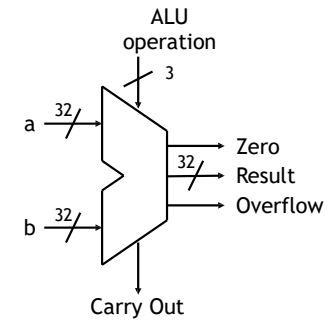
A whole MIPS ALU

f We can do the **subtraction** of $a - b$ by complementing the b input, setting **carry in** to 1, and performing an addition.

f The ALU also has a **slt** function, which sets **result** to 1 or 0 depending on whether or not a is less than b . This can be determined via subtraction.

f Finally, outputs can be generated to signal overflow or a result of zero.

ALU operation	Result
000	a and b
001	a or b
010	$a + b$
110	$a - b$
111	$a < b$



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MIPS ALU design

ALUs are a good example of modular hardware design.

- 32 full adders can be connected together to form a 32-bit adder.
- Bitwise operations can also be handled one bit at a time. Adders
- can also be used to perform subtraction.

Shift operations are usually handled outside the ALU by a combinational “barrel shifter” circuit, which can shift an arbitrary number of positions in constant time.

Summary

Signed numbers are represented using the **two's complement** format.

- A **sign bit** distinguishes between positive and negative numbers.
- Numbers are negated by complementing their bits and adding 1.
- This makes subtraction hardware easy to build.

Signed and unsigned numbers have many important differences.

- Some operations require that we **sign extend** the operands.
 - **Overflow** detection differs for signed and unsigned operations.
- The MIPS instruction set includes both signed and unsigned operations, which affect how constant operands and overflows are handled.