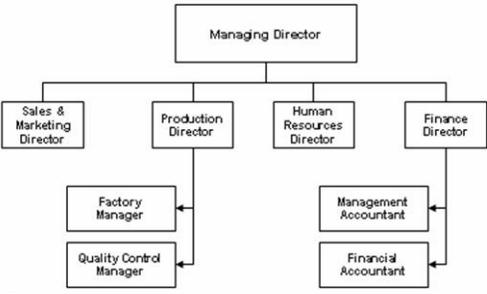
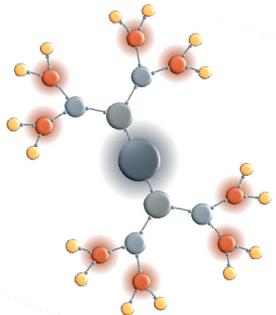




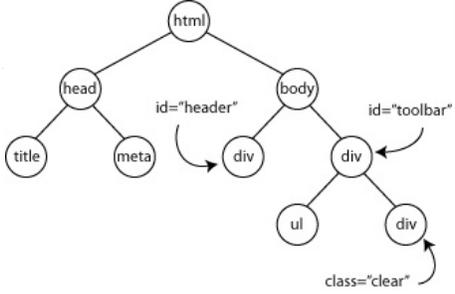
Binary Trees



Trees



```
graph TD; MD[Managing Director] --> SMD[Sales & Marketing Director]; MD --> PD[Production Director]; MD --> HRD[Human Resources Director]; MD --> FD[Finance Director]; PD --> FM[Factory Manager]; PD --> QCM[Quality Control Manager]; FD --> MA[Management Accountant]; FD --> FA[Financial Accountant];
```



```
graph TD; html((html)) --> head((head)); html --> body((body)); head --> title((title)); head --> meta((meta)); body --> div1((div)); body --> div2((div)); div1 --> header[id="header"]; div2 --> ul((ul)); div2 --> div3((div)); div3 --> clear[class="clear"]; div1 --> toolbar[id="toolbar"];
```



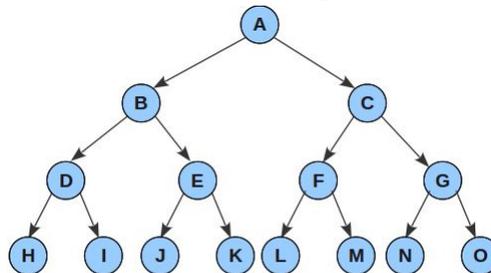
Advantages of Trees

- Trees reflect structural relationships in the data
- Trees are used to represent hierarchies
- Trees provide an efficient insertion and searching
- Trees are very flexible data structures, allowing to move subtrees around with minimum effort



A binary tree

- **Definition:** A binary tree is either **empty** or consists of a node called the **root** together with two binary trees called the **left subtree** and the **right subtree**.



Number of Leaves $(n) = 8$
Number of Nodes $(2n-1) = 2 \times 8 - 1 = 16 - 1 = 15$



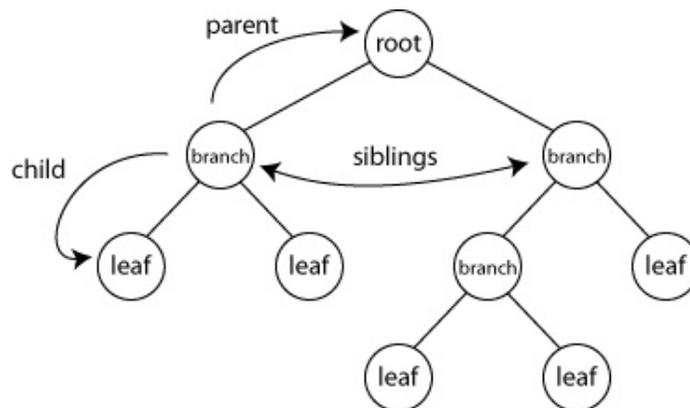
Parts of a binary tree

- A binary tree is composed of zero or more **nodes**
- Each node contains:
 - A **value** (some sort of data item)
 - A reference or pointer to a **left child** (may be **null**), and
 - A reference or pointer to a **right child** (may be **null**)
- A binary tree may be *empty* (contain no nodes)
- If not empty, a binary tree has a **root node**
 - Every node in the binary tree is reachable from the root node by a *unique* path
- A node with neither a left child nor a right child is called a **leaf**

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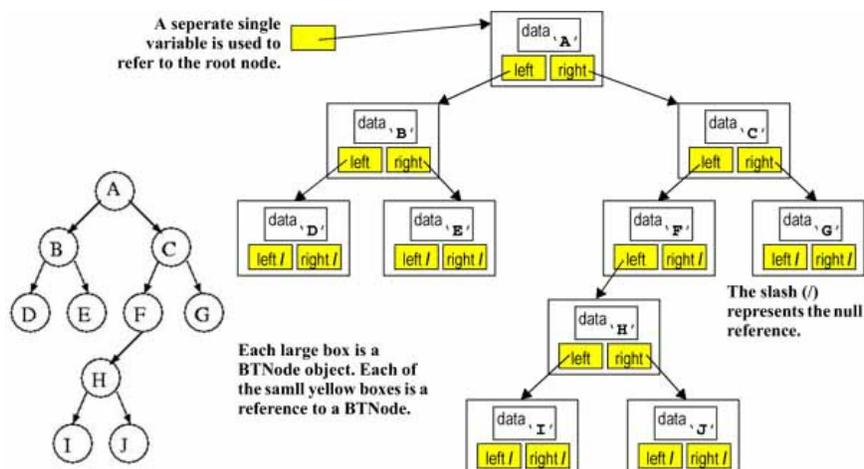


Parts of a binary tree





Picture of a binary tree

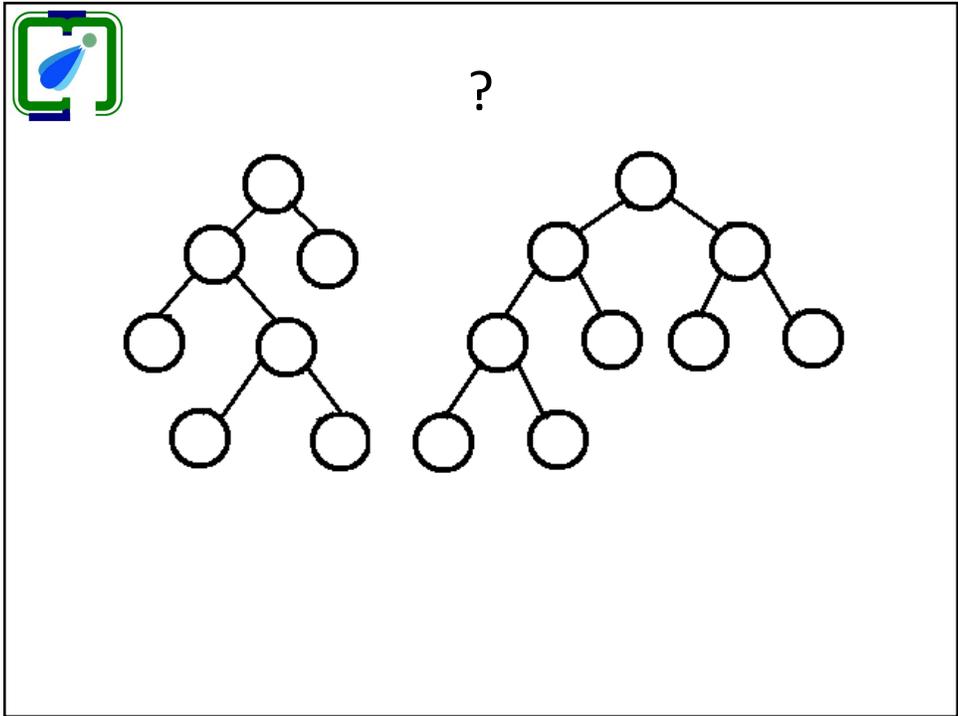


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Types of Binary Trees

- A **full binary tree** (sometimes **proper binary tree** or **2-tree** or **strictly binary tree**) is a tree in which every node other than the leaves has two children.
- A **perfect binary tree** is a *full binary tree* in which all *leaves* are at the same *depth* or same *level*, and in which every parent has two children
- A **complete binary tree** is a binary tree in which every level, *except possibly the last*, is completely filled, and all nodes are as far left as possible
- A **balanced binary tree** is commonly defined as a binary tree in which the *depth* of the two subtrees of every node never differ by more than 1




Size and depth (or height)

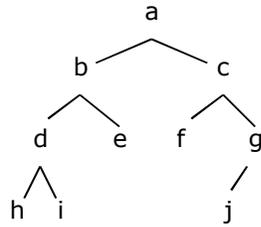
- The **size** of a binary tree is the number of nodes in it
 - This tree has size 12
- The **depth** of a node is its distance from the root
 - a is at depth zero
 - e is at depth 2
- The **depth** of a binary tree is the depth of its deepest node
 - This tree has depth 4

```

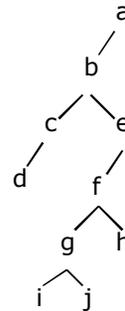
    graph TD
      a((a)) --- b((b))
      a --- c((c))
      b --- d((d))
      b --- e((e))
      c --- f((f))
      d --- g((g))
      e --- h((h))
      e --- i((i))
      f --- j((j))
      f --- k((k))
      i --- l((l))
    
```



Balance Binary Trees



A balanced binary tree



An unbalanced binary tree

- A binary tree is balanced if every level above the lowest is “full” (contains 2^n nodes)
- In most applications, a reasonably balanced binary tree is desirable

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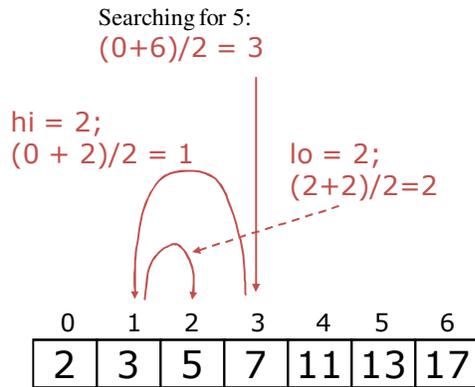
Height of a Binary Tree

- If h = height of a binary tree,
 max # of leaves = 2^h
 max # of nodes = $2^{h+1} - 1$

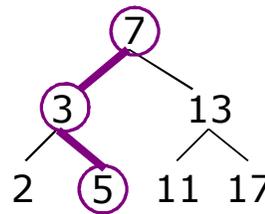


Binary search in an array

- Look at array location $(lo + hi)/2$



Using a binary search tree



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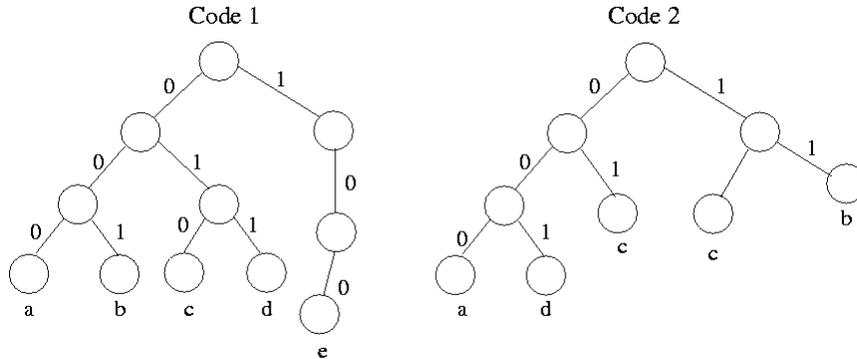
Huffman Coding – A Binary Tree Application [David A Huffman' 52]

- Suppose we have a message made from the five characters a, b, c, d, e, with probabilities 0.12, 0.40, 0.15, 0.08, 0.25, respectively

Symbol	Prob	code 1	code 2
a	0.12	000	000
b	0.40	001	11
c	0.15	010	01
d	0.08	011	001
e	0.25	100	10



Huffman's Coding cont...



Average length per character

$$\text{Code 1: } (0.12)(3) + (0.4)(3) + (0.15)(3) + (0.08)(3) + (0.25)(3) = 3$$

$$\text{Code 2: } (0.12)(3) + (0.4)(2) + (0.15)(2) + (0.08)(3) + (0.25)(2) = 2.2$$



Huffman's Coding cont...

- So the resulting “average length per character” = ?
- And the resulting Huffman tree for the example ?



Binary Tree traversals

- A binary tree is defined recursively: it consists of a **root**, a **left subtree**, and a **right subtree**
- To **traverse** (or **walk**) the binary tree is to visit each node in the binary tree exactly once
- Tree traversals are naturally recursive
- Since a binary tree has three “parts,” there are six possible ways to traverse the binary tree:
 - root, left, right
 - left, root, right
 - left, right, root
 - root, right, left
 - right, root, left
 - right, left, root

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Binary Tree Traversals

- Depth First Traversals
 - Preorder
 - Inorder
 - Postorder
- Breadth First Traversal



Preorder Traversal

- In **preorder**, the root is visited *first*
- Here's a preorder traversal to print out all the elements in the binary tree:

```
public void preorder(BinaryTree bt) {  
    if (bt == null) return;  
    printf(bt.value);  
    preorder (bt.leftChild);  
    preorder (bt.rightChild);  
}
```

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Inorder Traversal

- In **inorder**, the root is visited *in the middle*
- Here's an inorder traversal to print out all the elements in the binary tree:

```
public void inorder (BinaryTree bt) {  
    if (bt == null) return;  
    inorder(bt.leftChild);  
    printf(bt.value);  
    inorder(bt.rightChild);  
}
```

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Postorder Traversal

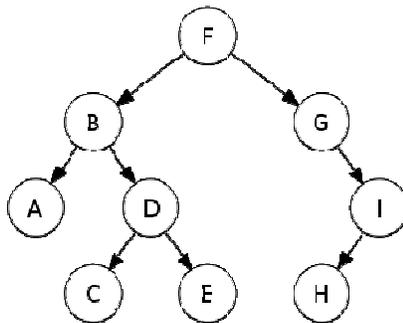
- In **postorder**, the root is visited *last*
- Here's a postorder traversal to print out all the elements in the binary tree:

```
public void postorder(BinaryTree bt) {
    if (bt == null) return;
    postorder(bt.leftChild);
    postorder(bt.rightChild);
    printf(bt.value);
}
```

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Tree Traversals – An Example



Depth-first

Preorder traversal sequence: F, B, A, D, C, E, G, I, H (root, left, right)

Inorder traversal sequence: A, B, C, D, E, F, G, H, I (left, root, right)

Postorder traversal sequence: A, C, E, D, B, H, I, G, F (left, right, root)

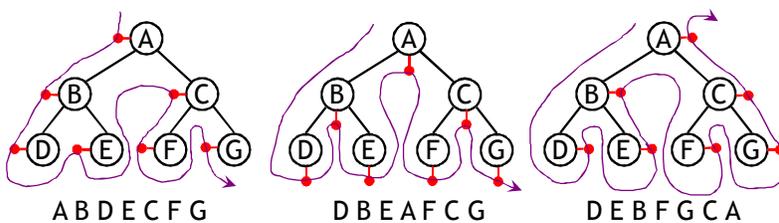


Tree traversals using “flags”

- The order in which the nodes are visited during a tree traversal can be easily determined by imagining there is a “flag” attached to each node, as follows:



- To traverse the tree, collect the flags:



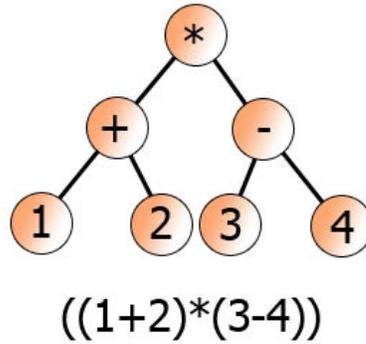
Other traversals

- The other traversals are the reverse of these three standard ones
 - That is, the right subtree is traversed before the left subtree is traversed
- Reverse preorder: root, right subtree, left subtree
- Reverse inorder: right subtree, root, left subtree
- Reverse postorder: right subtree, left subtree, root

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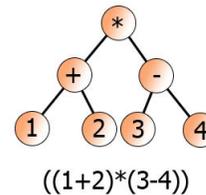


Binary Tree as Expression Tree



Binary Tree as Expression Tree

- Traversing the expression tree using
 - Preorder: results in to prefix code
 - In-order: results in to infix code (same expression)
 - Post-order: results in to postfix (reverse polish) code
- Find these codes for this example





Summary (Trees)

- A non-linear, hierarchical, and recursive data structures
- Form the basis of many useful and efficient data structures
- Traversals
 - Depth first
 - Pre-order
 - In-order
 - Post-order
 - Breadth First Traversal
- Applications of Binary Trees
 - Huffman coding
 - Expression Trees